

# PLO Compiler Example

February 22, 2021

# Example: Lexical Analysis

**Input** (sequence of characters):

**if**  $x < 0$  **then**  $z := -x$  **else**  $z := x$

**Output** (sequence of **lexical tokens**):

*KW\_IF, IDENTIFIER("x"), LESS, NUMBER(0), KW\_THEN,  
IDENTIFIER("z"), ASSIGN, MINUS, IDENTIFIER("x"), KW\_ELSE,  
IDENTIFIER("z"), ASSIGN, IDENTIFIER("x")*

# Example: Concrete syntax tree

```
if x < 0 then z := -x else z := x
```

The **grammar** of the language can be used to extract the **structure** of the program statement (from its sequence of lexical tokens).

The structure can be described using a **concrete syntax tree**.

if  $x < 0$  then  $z := -x$  else  $z := x$

Statement

Statement → *Assignment* | *CallStatement* | *ReadStatement* | *WriteStatement* |  
*WhileStatement* | *IfStatement* | *CompoundStatement*

...

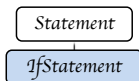
*IfStatement* → *KW\_IF* *Condition* *KW\_THEN* *Statement* *KW\_ELSE* *Statement*

*Condition* → *RelCondition*

*RelCondition* → *Exp* [ *RelOp* *Exp* ]

*Exp* → ...

if  $x < 0$  then  $z := -x$  else  $z := x$



*Statement* → *Assignment* | *CallStatement* | *ReadStatement* | *WriteStatement* |  
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...

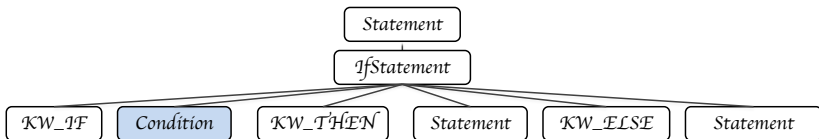
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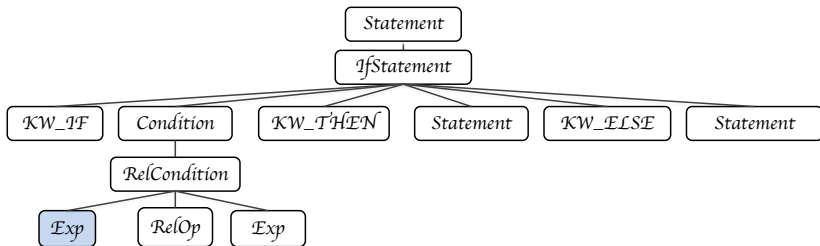
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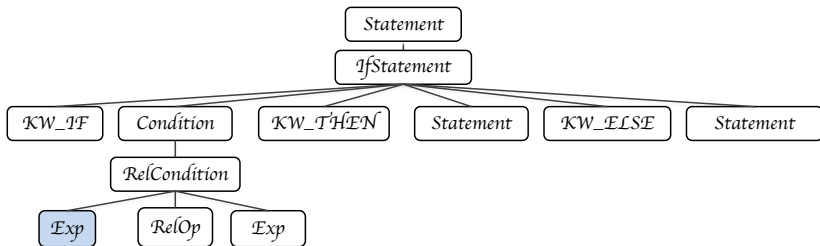
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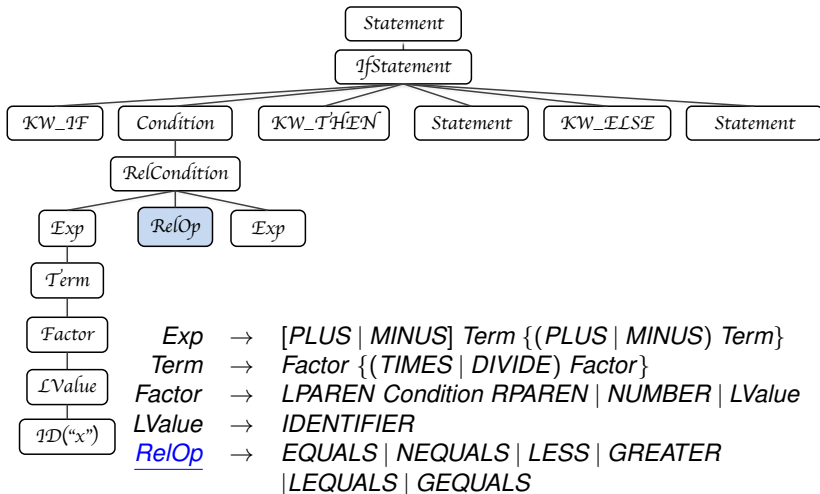
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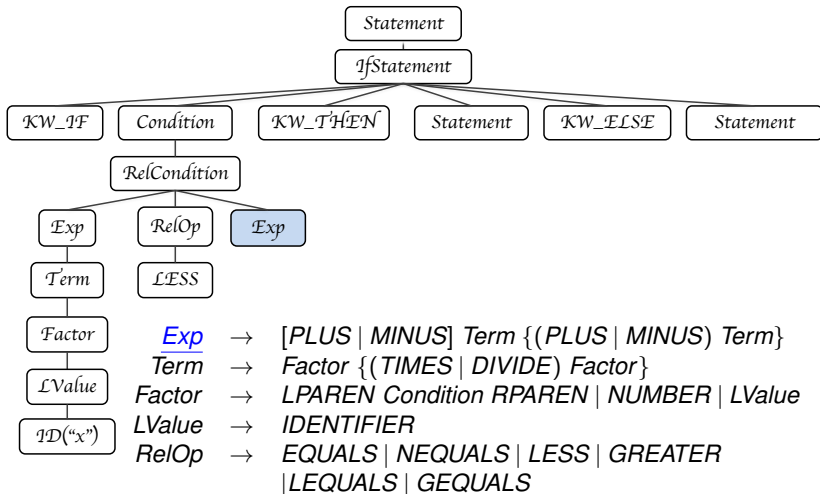
- $Exp$  →  $[PLUS \mid MINUS] Term \{(PLUS \mid MINUS) Term\}$   
 $Term$  →  $Factor \{(TIMES \mid DIVIDE) Factor\}$   
 $Factor$  →  $LPAREN Condition RPAREN \mid NUMBER \mid LValue$   
 $LValue$  →  $IDENTIFIER$   
 $RelOp$  →  $EQUALS \mid NEQUALS \mid LESS \mid GREATER$   
 $\mid LEQUALS \mid GEQUALS$



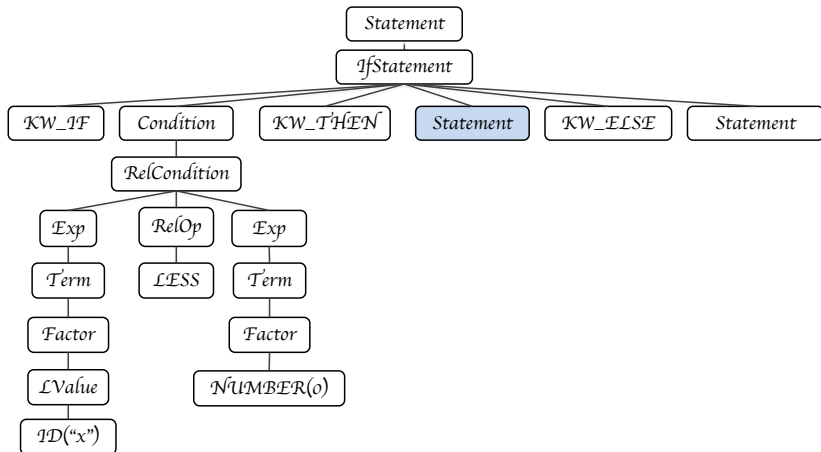
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**if**  $x < 0$  **then**  $z := -x$  **else**  $z := x$



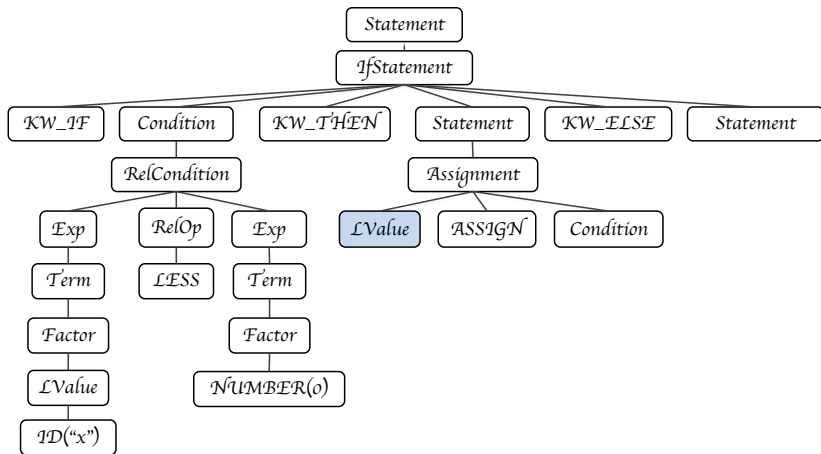
if  $x < 0$  then  $z := -x$  else  $z := x$



Statement → Assignment | CallStatement | ReadStatement | WriteStatement | WhileStatement | IfStatement | CompoundStatement

Assignment → LValue ASSIGN Condition

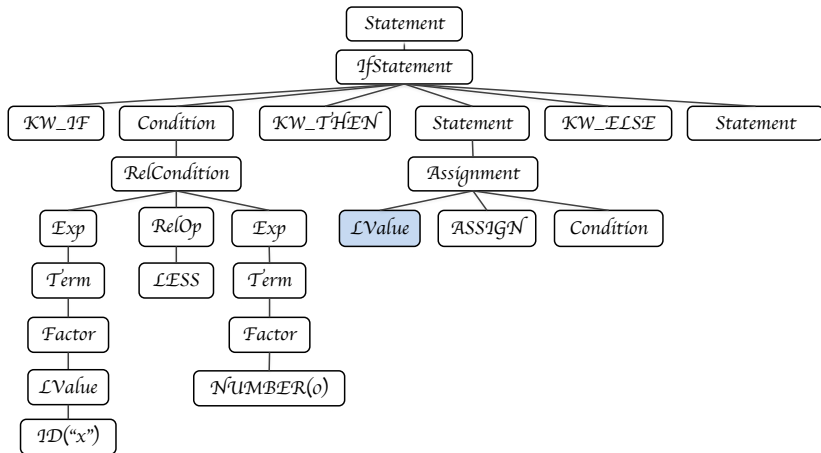
**if**  $x < 0$  **then**  $z := -x$  **else**  $z := x$



*Statement* → *Assignment* | *CallStatement* | *ReadStatement* | *WriteStatement* | *WhileStatement* | *IfStatement* | *CompoundStatement*

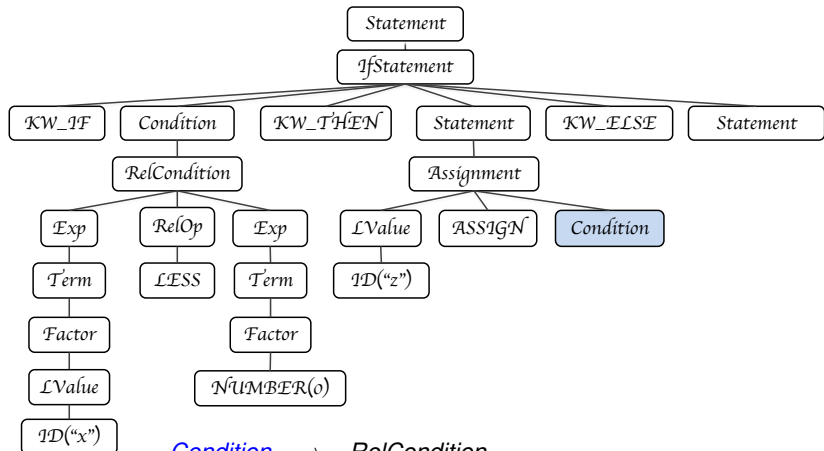
*Assignment* → *LValue* *ASSIGN* *Condition*

if  $x < 0$  then  $z := -x$  else  $z := x$



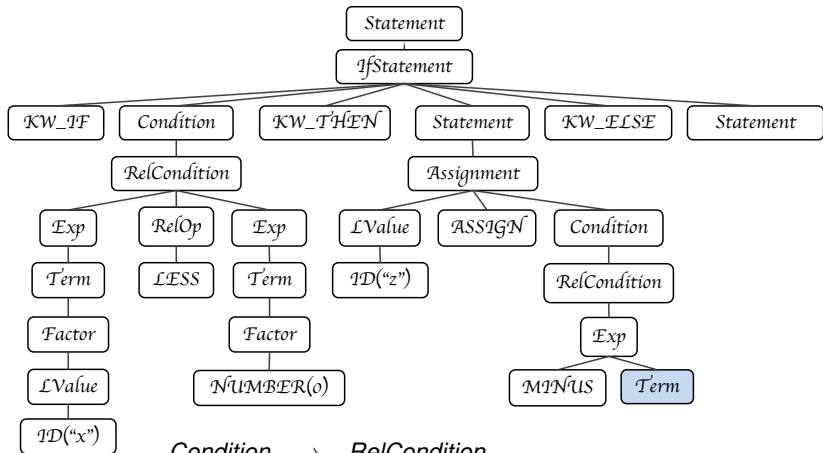
*LValue* → IDENTIFIER

if  $x < 0$  then  $z := \underline{-x}$  else  $z := x$



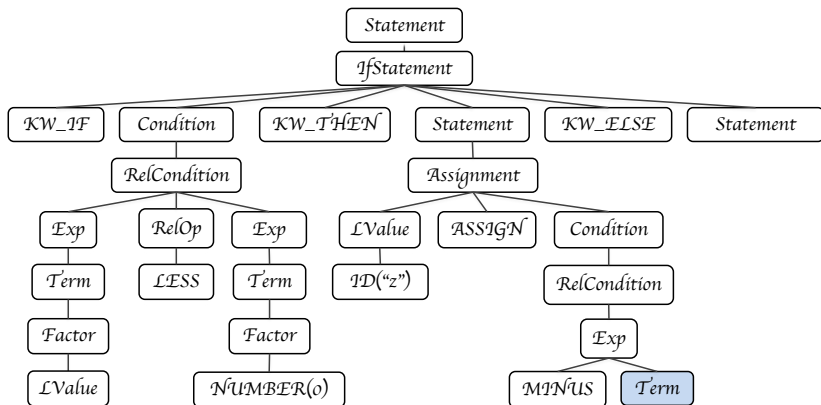
- Condition → *RelCondition*
- RelCondition* → *Exp [ RelOp Exp ]*
- Exp* → *[ PLUS | MINUS ] Term { ( PLUS | MINUS ) Term }*
- Term* → ...

if  $x < 0$  then  $z := -x$  else  $z := x$



*Condition* → *RelCondition*  
*RelCondition* → *Exp* [ *RelOp* *Exp* ]  
*Exp* → [ *PLUS* | *MINUS* ] *Term* { ( *PLUS* | *MINUS* ) *Term* }  
*Term* → ...

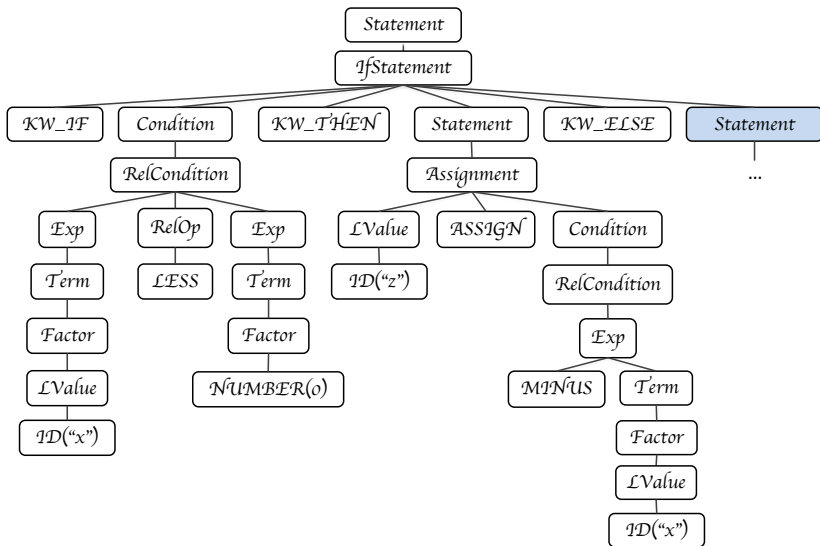
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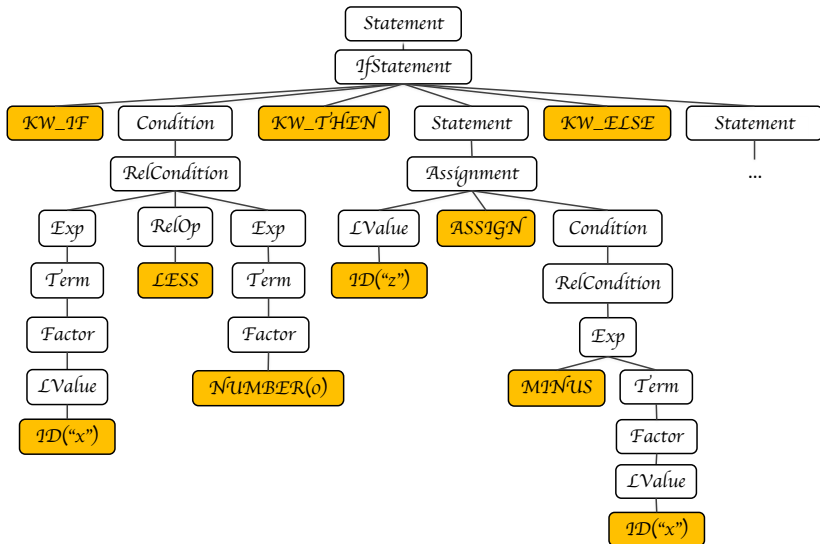
- Term** → *Factor* { (TIMES | DIVIDE) *Factor* }
- Factor** → LPAREN *Condition* RPAREN | NUMBER | LValue
- LValue** → IDENTIFIER



**if  $x < 0$  then  $z := -x$  else  $z := x$**

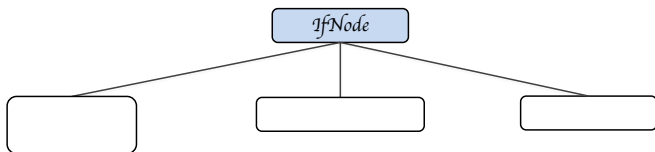


**if  $x < 0$  then  $z := -x$  else  $z := x$**



# Example: Abstract Syntax Tree

if  $x < 0$  then  $z := -x$  else  $z := x$



StatementNode(Location loc) with subclasses

AssignmentNode(ExpNode lvalue, ExpNode e)

IfNode(ExpNode cond, StatementNode s1, StatementNode s2)

...

ExpNode(Location loc, Type t) with subclasses

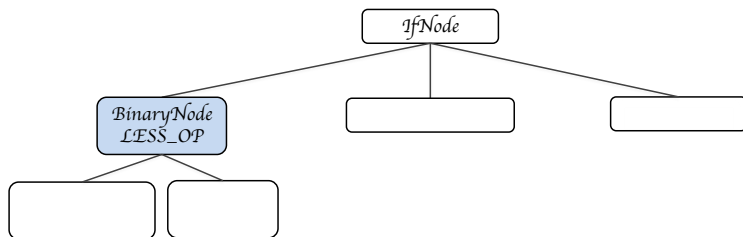
ConstNode(int value)

IdentifierNode(String id)

BinaryNode(Operator op, ExpNode left, ExpNode right)

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if  $x < 0$  then  $z := -x$  else  $z := x$



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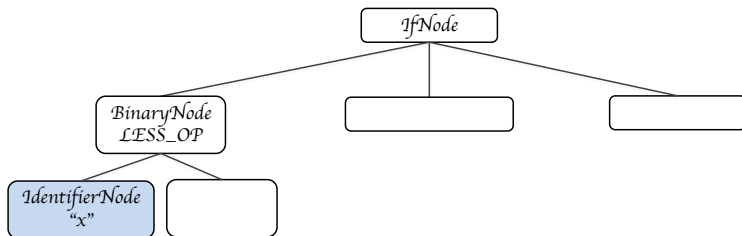
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ExpNode(Location loc, Type t) with subclasses

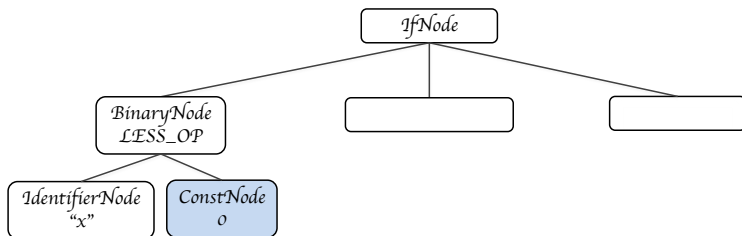
ConstNode(int value)

[IdentifierNode](#)(String id)

BinaryNode(Operator op, [ExpNode left](#), ExpNode right)

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if  $x < 0$  then  $z := -x$  else  $z := x$



StatementNode(Location loc) with subclasses

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ExpNode(Location loc, Type t) with subclasses

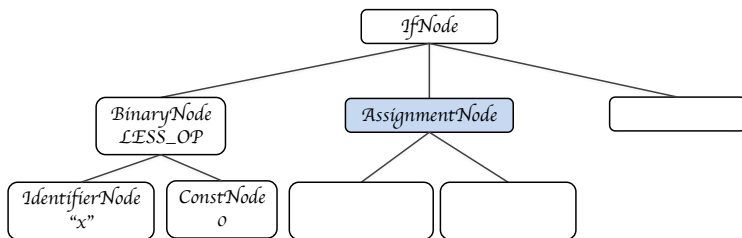
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StatementNode(Location loc) with subclasses

[AssignmentNode](#)(ExpNode lvalue, ExpNode e)

[IfNode](#)(ExpNode cond, [StatementNode s1](#), StatementNode s2)

...

ExpNode(Location loc, Type t) with subclasses

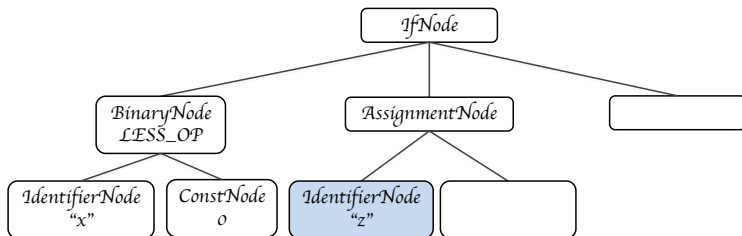
ConstNode(int value)

IdentifierNode(String id)

UnaryNode(Operator op, ExpNode arg)

# Example: Abstract Syntax Tree

if  $x < 0$  then  $z := -x$  else  $z := x$



StatementNode(Location loc) with subclasses

AssignmentNode([ExpNode lvalue](#), ExpNode e)

IfNode(ExpNode cond, StatementNode s1, StatementNode s2)

...

ExpNode(Location loc, Type t) with subclasses

ConstNode(int value)

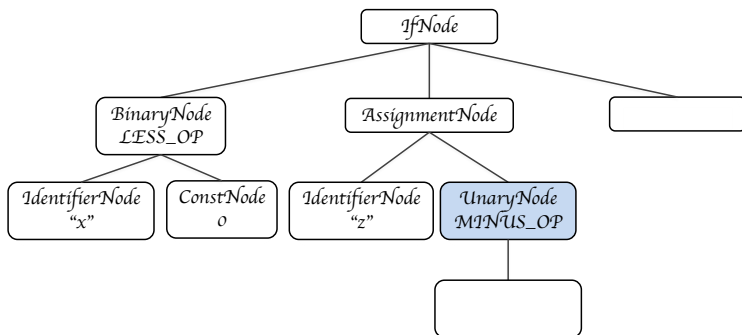
[IdentifierNode](#)(String id)

UnaryNode(Operator op, ExpNode arg)



# Example: Abstract Syntax Tree

if  $x < 0$  then  $z := \underline{-x}$  else  $z := x$



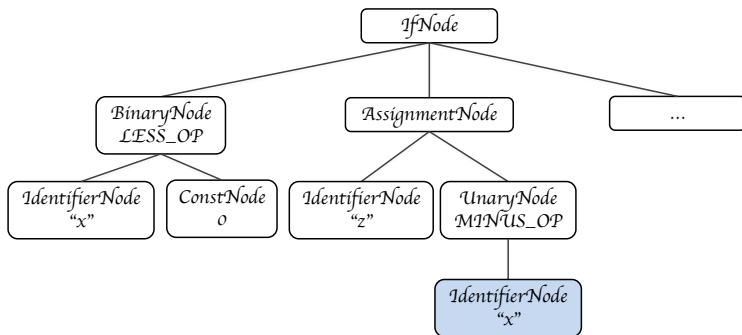
StatementNode(Location loc) with subclasses  
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...

ExpNode(Location loc, Type t) with subclasses  
[UnaryNode](#)(Operator op, ExpNode arg)

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ExpNode(Location loc, Type t) with subclasses

ConstNode(int value)

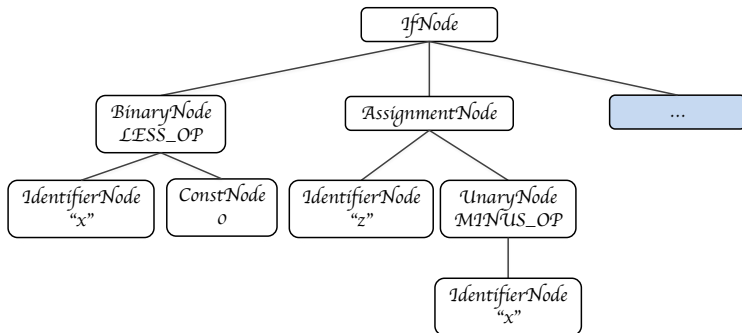
[IdentifierNode](#)(String id)

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# Example: Abstract Syntax Tree

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StatementNode(Location loc) with subclasses

[AssignmentNode](#)(ExpNode lvalue, ExpNode e)

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...

# Type Checking the AST

[IdentifierNode](#) is used during parsing to represent a reference to either a symbolic constant or a variable.

As part of the static semantics (type checking) phase it will be transformed to either a [ConstNode](#) or a [VariableNode](#).

```
ExpNode(Location loc, Type t) with subclasses
  ConstNode(int value)
  IdentifierNode(String id)
  VariableNode(SymEntry.VarEntry entry)
  ...
```

# Type Checking the AST

A number of other node types are only introduced in the static semantics phase:

- a [DereferenceNode](#) represents a dereference of a variable address (left value) to access its (right) value;
- an expression of a type,  $T$ , can be narrowed to a subrange of  $T$  ([NarrowSubrangeNode](#)).
- an expression of a subrange type can be widened to the base type of the subrange ([WidenSubrangeNode](#));

ExpNode(Location loc, Type t) with subclasses

DereferenceNode(ExpNode leftValue)

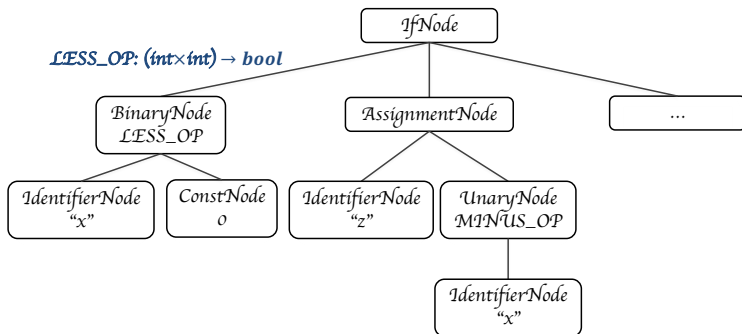
NarrowSubrangeNode(ExpNode e)

WidenSubrangeNode(ExpNode e)

...

# Example: Type Checking the AST

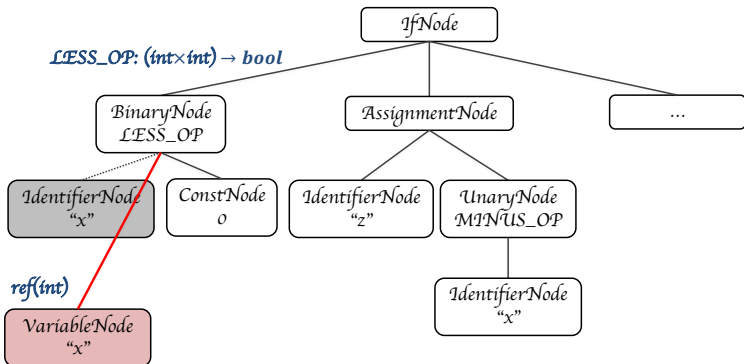
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ExpNode(Location loc, Type t) with subclasses  
 VariableNode(SymEntry.VarEntry entry)  
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 ...

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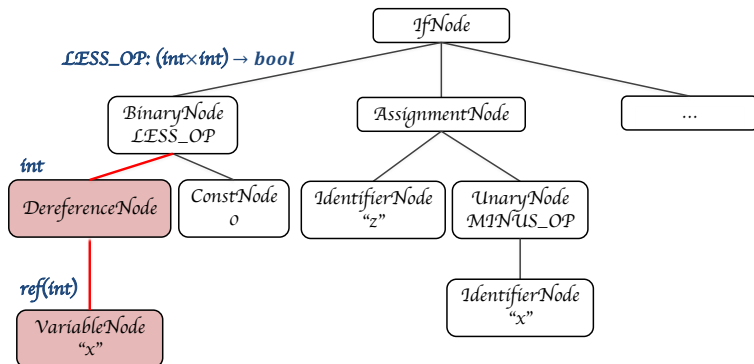
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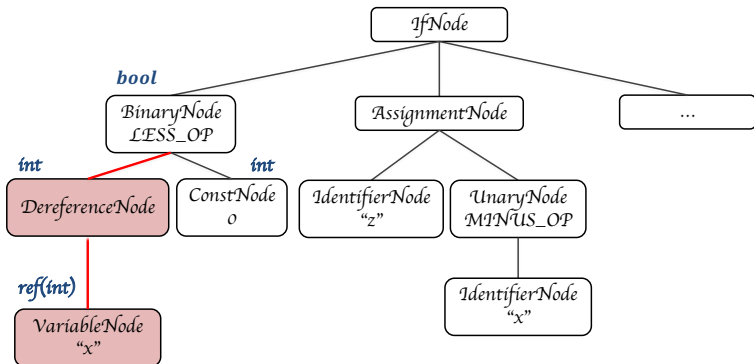


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# Example: Type Checking the AST

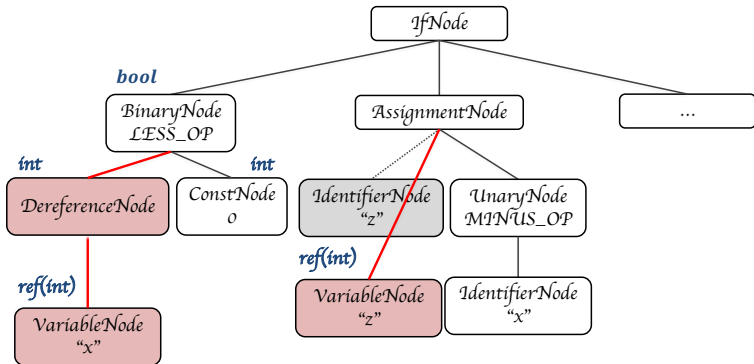
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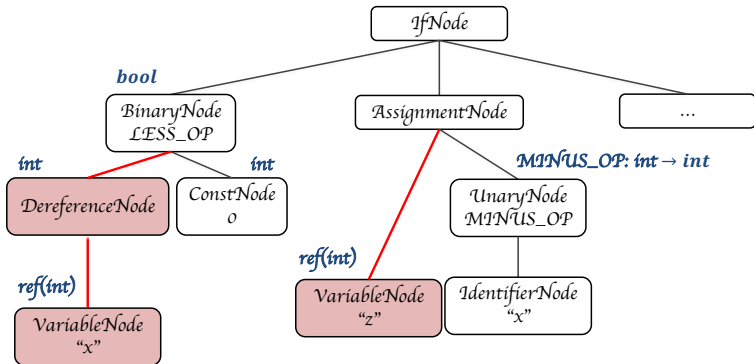


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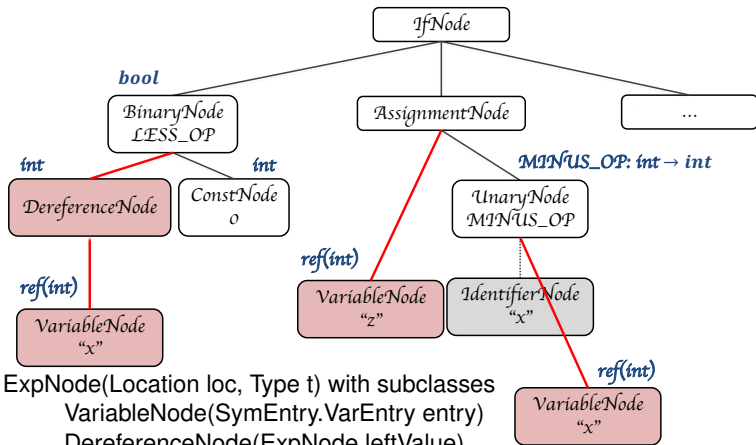


ExpNode(Location loc, Type t) with subclasses  
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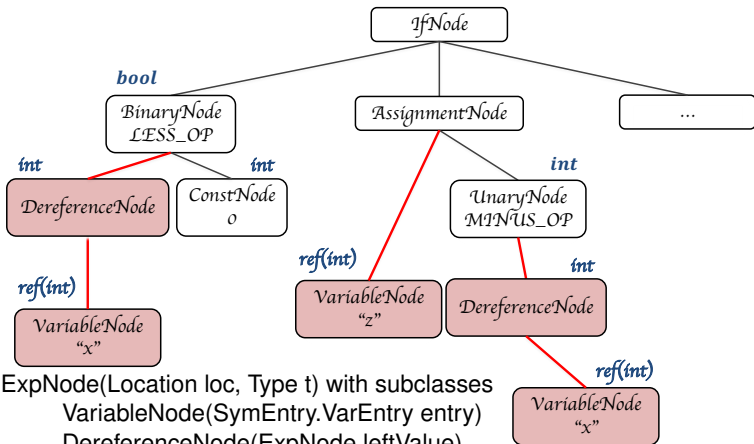
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# Example summary

**if**  $x < 0$  **then**  $z := -x$  **else**  $z := x$

- Lexical Analysis of the sequence of program characters produced a sequence of lexical tokens.
- Syntax Analysis of the sequence of lexical tokens was used to produce:
  - A concrete syntax tree – **not produced by the compiler!**
  - An abstract syntax tree (AST)
- Static Analysis was used to:
  - resolve references to identifiers;
  - type check the abstract syntax tree; and
  - update the abstract syntax tree with type coercions.

# Lexical Analysis

**Input** (sequence of characters):

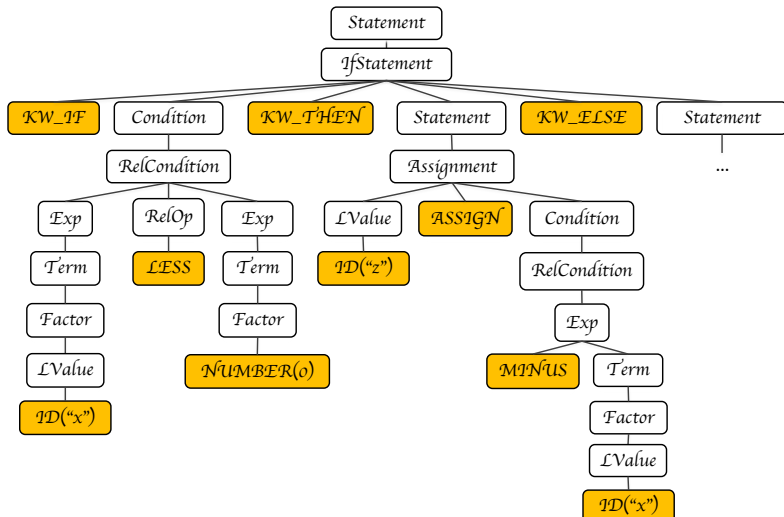
```
if x < 0 then z := -x else z := x
```

**Output** (sequence of **lexical tokens**):

```
KW_IF, IDENTIFIER("x"), LESS, NUMBER(0), KW_THEN,  
IDENTIFIER("z"), ASSIGN, MINUS, IDENTIFIER("x"), KW_ELSE,  
IDENTIFIER("z"), ASSIGN, IDENTIFIER("x")
```

# Syntax Analysis: Concrete Syntax Tree

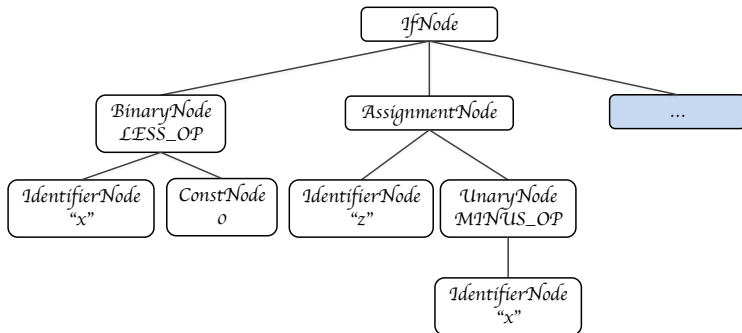
**if  $x < 0$  then  $z := -x$  else  $z := x$**





# Syntax Analysis: Abstract Syntax Tree

**if**  $x < 0$  **then**  $z := -x$  **else**  $z := x$



# Static Analysis: Type Checking the AST

**if  $x < 0$  then  $z := -x$  else  $z := x$**

